



FIGURE 1

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contend():
s ← 0;
repeat [
    Q1 ← ∅;
    ||u ∈ U [ statusu ← u, try();
        Q1 ← {u} ∪ Q1;
    ] until (∃ Q ∈ Q: Q ⊆ Q1);
    if (|{u : statusu = LOCKED}| ≤ b)
        return;
    else
        s ← s + 1;
        d ← R[(Δ+4δ)]... 2s(Δ+4δ);
        sleep(d);
] until (false);

```

FIGURE 2A: MUTUAL EXCLUSION PROTOCOL CLIENT PROGRAM

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try():
    if (clock() - lastGranted >  $\Delta + 2 \delta$ )
        lastGranted  $\leftarrow$  clock();
        return FREE;
    else
        return LOCKED;

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FIGURE 2B: MUTUAL EXCLUSION PROTOCOL SERVER PROGRAM

<pre> 1) submit(o): 2) waiting ← true; 3) $\parallel V_1 \leftarrow \emptyset; Q_1 \leftarrow_R Q;$ 4) $\parallel_{u \in Q_1} [p_u \leftarrow u.submit(o); V_1 \leftarrow \{p_u\} \cup V_1;$ 5) $] \text{ until } (\exists p: \{p_u \in V_1 : p=p_u\} \geq b+1;$ 6) $\text{ waiting} \leftarrow \text{false};$ 7) $\text{ return } p: \{p_u \in V_1 : p=p_u\} \geq b+1$ 8) $\parallel \text{ repeat } [$ 9) $\text{ contend}();$ 10) $Q_2 \leftarrow \emptyset;$ 11) $\parallel_{u \in U} [\langle \sigma_u^c, \sigma_u^{pc}, \text{proposer}_u, \text{pending}_u \rangle \leftarrow u.get(r);$ 12) $Q_2 \leftarrow \{u\} \cup Q_2;$ 13) $] \text{ until } (\exists Q \in Q: Q \subseteq Q_2);$ 14) $\sum^c \leftarrow \{\sigma' : \{u: \sigma' = \sigma_u^c\} \geq b+1\};$ 15) $\sigma^c \leftarrow \sigma: \sigma.version = \max_{\sigma' \in \sum^c} \{\sigma'.version\};$ 16) $\sigma^{pc} \leftarrow \text{choose}(\{\langle \sigma_u^{pc}, \text{proposer}_u \rangle: \sigma_u^{pc} \neq \perp\});$ 17) $\text{completed} \leftarrow \max\{\text{completed}, \max\{v: \{u: \sigma_u^{pc}.version > v\} \geq b+1\}\};$ 18) $\text{if } \sigma^c \neq \perp \wedge \sigma^c.version > \text{completed}$ 19) $\quad \sigma \leftarrow \sigma^c$ 20) $\text{else if } (\sigma^{pc} \neq \perp \wedge \sigma^{pc}.version > \text{completed})$ 21) $\quad \sigma \leftarrow \sigma^{pc};$ 22) else 23) $\quad \text{pending} \leftarrow \{o: \exists \{u: o \in \text{pending}_u\} \geq b+1\};$ 24) $\quad \sigma \leftarrow \text{apply}(\text{pending}, \sigma^c);$ 25) $Q_2 \leftarrow \emptyset;$ 26) $\parallel_{u \in U} [u.propose(\sigma', r); Q_2 \leftarrow \{u\} \cup Q_2;$ 27) $] \text{ until } (\exists Q \in Q: Q \subseteq Q_2);$ 28) $Q_2 \leftarrow \emptyset;$ 29) $\parallel_{u \in U} [u.commit(\sigma, r); Q_2 \leftarrow \{u\} \cup Q_2;$ 30) $] \text{ until } (\exists Q \in Q: Q \subseteq Q_2);$ 31) $\text{completed} \leftarrow \max\{\text{completed}, \sigma.version\};$ 32) $] \text{ until } (\text{waiting} = \text{false});$ </pre>	<pre> 33) $\text{choose}(\{\langle \sigma_u^{pc}, \text{proposer}_u \rangle\} u):$ 34) $S[1, 2, \dots] \leftarrow \{\langle \sigma_u^{pc}, \text{proposer}_u \rangle\}_u$ sorted in descending order by $\langle \sigma, \text{proposer} \rangle > \langle \sigma', \text{proposer}' \rangle \Leftrightarrow (\sigma.version > \sigma'.version \vee (\sigma.version = \sigma'.version \wedge \text{proposer} > \text{proposer}'))$; 35) $i \leftarrow 1; \text{count} \leftarrow [0, 0, \dots];$ 36) repeat 37) $\quad [\langle \sigma, \text{proposer} \rangle \leftarrow S[i];$ 38) $\quad \text{count}[\sigma] \leftarrow \text{count}[\sigma] + 1;$ 39) $\quad i \leftarrow i + 1;$ 40) $\quad] \text{ until } (\exists \sigma: \text{count}[\sigma] \geq b+1 \vee i > S);$ 41) $\quad \text{if } (\exists \sigma: \text{count}[\sigma] \geq b+1)$ 42) $\quad \quad \text{return } \sigma: \text{count}[\sigma] \geq b+1;$ 43) $\quad \text{else}$ 44) $\quad \quad \text{return } \perp;$ 45) $\text{apply}(\text{pending}, \sigma):$ 46) $\text{repeat } [o \leftarrow_R \text{pending};$ 47) $\quad \text{pending} \leftarrow \text{pending} \setminus \{o\};$ 48) $\quad \text{if } (\sigma'.reflects(o) = \text{false})$ 49) $\quad \quad \sigma'.doOp(o);$ 50) $\quad] \text{ until } (\text{pending} = \emptyset);$ 51) $\sigma'.version \leftarrow \sigma'.version + 1;$ 52) $\text{return } \sigma';$ </pre>
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FIGURE 3A: CLIENT SIDE OF AN ORDERING PROTOCOL

<pre> 1) submit (o): 2) pending ← pending ∪ {o}; 3) sleep until (response(o) ≠ ⊥); 4) return response (o); </pre>	<pre> 5) get(r): 6) if @ > maxRank) 7) maxRank ← r; 8) return ⟨∅, σ^{pc}, σ^{pc}, prosper, pending⟩ 9) else 10) throw RankException; </pre>
<pre> 11) propose (σ, r): 12) if @ ≥ maxRank) 13) maxRank ← r; 14) proposer ← r; 15) σ^{pc} ← σ; 16) return; 17) else 18) throw RankException; </pre>	<pre> 19) commit (σ, r): 20) if @ ≥ maxRank) 21) maxRank ← r; 22) σ^c, σ; 23) pending ← pending \ {o:σ.reflects(o)=true}; 24) response ← response σ. response; 25) return; 26) else 27) throw RankException; </pre>

FIGURE 3B: SERVER SIDE OF AN ORDERING PROTOCOL